

## Evolution of Binding Sites

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Molecular Information Theory Group

132 p53 binding sites



El Duomo, Florence, Italy


## number of number of symbols bits example

M B

2
1


4
2


83

$M=2^{B} \quad B=\log _{2} M$


## Information Theory: One-Minute Lesson

number of number of symbols bits example

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number of number of symbols B

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1

4

| $M$ | $B$ |
| :--- | :--- |
| 2 | 1 |
| 4 | 2 |
| 8 | 3 |
| $M=2^{B}$ | $B=\log _{2} M$ |



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$M=2^{B} \quad B=\log _{2} M$


## Sequence Logo

Bacteriophage T7 RNA polymerase binding sites


1 ttattaatacaactcactataaggagag
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3 cggttaatacgactcactataggagaac
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6 of 17 sites
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More Information Theory - 1

## An Intuitive Approach

Information to chose one symbol from $M$ symbols:

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\log _{2} M \tag{1}
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$1 / M$ is like the probability of a symbol.

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$$

$1 / M$ is like the probability of a symbol.
If the probabilities $P_{i}$ of different symbols, $i$, are not equal, then the surprisal is:

$$
\begin{equation*}
u_{i} \equiv-\log _{2} P_{i} . \tag{2}
\end{equation*}
$$

how surprised one is to see a symbol

## EXAMPLE

A phone rings once every 1024 seconds.


$$
\begin{align*}
P_{\text {ring }} & =1 / 1024  \tag{3}\\
P_{\text {silent }} & =1023 / 1024 \tag{4}
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## More Information Theory - 2

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Surprisal:

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The average surprisal is called the uncertainty, $H$ :

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H=P_{\text {ring }} \times \text { surprisal }_{\text {ring }}
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\end{gather*}
$$

For $M$ symbols use the sum $\left(\sum\right)$ notation:

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H=\sum_{i=1}^{M} P_{i} \times\left(\text { surprisal for } P_{i}\right)
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& =-\sum_{i=1}^{M} P_{i} \log _{2} P_{i} \quad \text { bits per symbol }
\end{align*}
$$

## More Information Theory - 4

Information is a decrease in uncertainty

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R=H_{\text {before }}-H_{\text {after }} \tag{12}
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132 p53 binding sites


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Note: with only one base, $H_{\text {after }}=0$ so $R=2$ bits/base.

## Information required to find a set of binding sites

$G=\#$ of potential binding sites

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R_{\text {frequency }} & =H_{\text {before }}-H_{a f t e r} \\
& =\log _{2} G-\log _{2} \gamma \\
& =-\log _{2} \gamma / G
\end{aligned}
$$

## Rfrequency

Information required to find a set of binding sites in a genome


16 positions
1 site
$\log _{2} 16 / 1=4$ bits


16 positions
2 sites
$\log _{2} 16 / 2=3$ bits

Donor and acceptor logos

donor


## Rsequence and Rfrequency for Splice Acceptors

$R_{\text {sequence }}$

- Information at binding site sequences (area under sequence logo)
- from: binding site sequences
- 9.4 bits per site


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## $R_{\text {sequence }}$



- Information at binding site sequences (area under sequence logo)
- from: binding site sequences
- 9.4 bits per site
$R_{\text {frequency }}$

- Information needed to locate the sites
- from: size of genome and number of sites (length of intron+exon)
- 9.7 bits per site

$$
R_{\text {frequency }} / R_{\text {sequence }}=0.97
$$

Rsequence $=$ Rfrequency Hypothesis

> Hypothesis:
> The information in binding site patterns is just sufficient for the sites to be found in the genome

| Binding Site <br> Recognizer |  |  |  |
| :--- | :--- | :---: | :--- |
|  | Total Pattern <br> Information <br> $=\mathbf{R}_{\text {sequence }}$ <br> (bits) | Information needed to <br> Locate <br> $=\mathbf{R}_{\text {frequency }}$ <br> (bits) | $\frac{\text { Pattern Info }}{\text { Location Info }}$ <br> $=\frac{\mathbf{R}_{\text {sequence }}}{\mathbf{R}_{\text {frequency }}}$ |
| Spliceosome acceptor $^{2}$ | $\mathbf{9 . 3 5} \pm \mathbf{0 . 1 2}$ | $\mathbf{9 . 6 6}$ | $\mathbf{0 . 9 7} \pm \mathbf{0 . 0 1}$ |
| Spliceosome donor | $\mathbf{7 . 9 2} \pm \mathbf{0 . 0 9}$ | $\mathbf{9 . 6 6}$ | $\mathbf{0 . 8 2} \pm \mathbf{0 . 0 1}$ |
| Ribosome | $\mathbf{1 1 . 0}$ | $\mathbf{1 0 . 6}$ | $\mathbf{1 . 0}$ |
| $\lambda$ cl/cro | $\mathbf{1 7 . 7} \pm \mathbf{1 . 6}$ | $\mathbf{1 9 . 3}$ | $\mathbf{0 . 9} \pm \mathbf{0 . 1}$ |
| LexA | $\mathbf{2 1 . 5} \pm \mathbf{1 . 7}$ | $\mathbf{1 8 . 4}$ | $\mathbf{1 . 2} \pm \mathbf{0 . 1}$ |
| TrpR | $\mathbf{2 3 . 4} \pm \mathbf{1 . 9}$ | $\mathbf{2 0 . 3}$ | $\mathbf{1 . 2} \pm \mathbf{0 . 1}$ |
| Lacl | $\mathbf{1 9 . 2} \pm \mathbf{2 . 8}$ | $\mathbf{2 1 . 9}$ | $\mathbf{0 . 9} \pm \mathbf{0 . 1}$ |
| ArgR | $\mathbf{1 6 . 4}$ | $\mathbf{1 8 . 4}$ | $\mathbf{0 . 9}$ |
| O $(\lambda$ Origin) | $\mathbf{2 0 . 9}$ | $\mathbf{1 9 . 9}$ | $\mathbf{1 . 0}$ |
| AraC | $\mathbf{1 9 . 3}$ | $\mathbf{1 9 . 3}$ | $\mathbf{1 . 0}$ |
| Transcription at TATA ${ }^{3}$ | $\mathbf{3 . 3}$ | $\sim \mathbf{3}$ | $\sim \mathbf{1}$ |
| T7 Promoter | $\mathbf{3 5 . 4}$ | $\mathbf{1 6 . 5}$ | $\mathbf{2 . 1}$ |

[^0]
## $R_{\text {sequence }}$ versus $R_{\text {frequency }}$-meaning

The information in the binding site pattern ( $R_{\text {sequence }}$ ) is close to
The information needed to find the binding sites ( $R_{\text {frequency }}$ )

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But for a species in a stable environment:

- size of genome $(G)$ is fixed (e. g. E. coli has $4.7 \times 10^{6}$ bp)
- number of binding sites $(\gamma)$ is fixed (e. g. there are $\sim 50$ E. coli LexA sites) so $R_{\text {frequency }}=\log _{2} G / \gamma$ is fixed


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Rsequence must evolve towards Rfrequency!

## Evolution of Binding Sites

- $\quad R_{\text {frequency }}$ is fixed relative to $R_{\text {sequence }}$


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- $\quad R_{\text {frequency }}$ is fixed relative to $R_{\text {sequence }}$
- Does $R_{\text {sequence }}$ evolve toward $R_{\text {frequency }}$ ?
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Setup a Computer Model, 'Ev':
A population of "creatures" with

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- genomes containing 4 bases (A, C, G, T)
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- genomes containing 4 bases (A, C, G, T)
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$R_{\text {frequency }}$ is fixed
- a recognizer gene encoded in the sequence: use a weight matrix


## How A Weight Matrix Works

Sequence matrix, $s(b, l, j)$ for sequence $j$

| base b | position 1 |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | C | A | G | G | T | C | T | G | C | A |
|  | -3 | -2 | -1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
| A | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| C | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| G | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 |
| T | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 |

Individual information weight matrix, $R_{i w}(b, l)$

| base b | position 1 |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | -3 | -2 | -1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
| A | +0.4 | +1.3 | -1.4 | -8.8 | -5.8 | +1.1 | +1.5 | -1.8 | -0.7 | +0.0 |
| C | +0.6 | -0.8 | -2.4 | -7.8 | -5.5 | -3.7 | -1.6 | -2.2 | -0.5 | -0.2 |
| G | -0.6 | -1.0 | +1.6 | $\boxed{+2.0}$ | -6.2 | +0.7 | -1.1 | +1.7 | -0.3 | +0.4 |
| T | -1.0 | -0.9 | -1.7 | -5.8 | +2.0 | -3.4 | -1.6 | -2.2 | +0.9 | -0.5 |

## How A Weight Matrix Works

Sequence matrix, $s(b, l, j)$ for sequence $j$


Individual information weight matrix, $R_{i w}(b, l)$

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| G | -0.6 | -1.0 | +1.6 | -2.0 | -6.2 | +0.7 | -1.1 | +1.7 | -0.3 | +0.4 |
| T | -1.0 | -0.9 | -1.7 | -5.8 | $\boxed{+2.0}$ | -3.4 | -1.6 | -2.2 | +0.9 | -0.5 |



## Sequence Walker



## Unevolved Ev Creature



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Genome positions available $G=256$ bases $R_{\text {frequency }}=\log _{2} 256 / 16=4$ bits

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## Evolution Cycle

- EVALUATE each creature
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$\square$ finding a site at a wrong place



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$\square$ missing a site at a right place
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- Sort the creatures by their mistakes



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- REPLICATE: the best creatures are duplicated and replace the worst ones



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- REPLICATE: the best creatures are duplicated and replace the worst ones
- MUTATE all genomes randomly









```
C A A T T T G T A C A A A C T G A AlGA C A G G
```





Evolution of Binding Sites


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## Mathematics of Evolution 1

## Shannon Information Measure of Binding Site Patterns

Information is measured as a decrease in uncertainty:

$$
R=H_{\text {before }}-H_{\text {after }} \quad \text { (bits per symbol) }
$$

Before binding there are 4 possible bases at each position $l$, so the uncertainty is:

$$
\begin{aligned}
H_{\text {before }}(l) & =\log _{2} 4 \quad(\text { bits per base }) \quad(16) \\
& \approx 2
\end{aligned}
$$

## Mathematics of Evolution 2

Before binding there are 4 possible bases at each position $l$, so the uncertainty is:

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& \approx 2
\end{aligned}
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After binding the uncertainty depends on the frequencies of bases $b$ at positions $l$ in a binding site, $f(b, l)$ :

$$
\begin{aligned}
H_{a f t e r}(l)= & -\sum_{b \in\{A, C, G, T\}} f(b, l) \log _{2} f(b, l) \quad(17) \\
& (\text { bits per base })
\end{aligned}
$$

## The information at a position $l$ is:

$$
\begin{aligned}
R_{\text {sequence }}(l) & =H_{\text {before }}(l)-H_{\text {after }}(l) \\
& \approx 2-H_{\text {after }}(l) \quad \text { (bits per base) }
\end{aligned}
$$

## Mathematics of Evolution 3

## The information at a position $l$ is:

$$
\begin{align*}
R_{\text {sequence }}(l) & =H_{\text {before }}(l)-H_{\text {after }}(l)  \tag{18}\\
& \approx 2-H_{\text {after }}(l) \quad \text { (bits per base) }
\end{align*}
$$

The total site information is:

$$
\begin{aligned}
R_{\text {sequence }} & =\sum_{l}\left(H_{\text {before }}(l)-H_{\text {after }}(l)\right) \\
& \approx 2 l-H_{\text {after }} \quad(\text { bits per site })(19)
\end{aligned}
$$

During evolution, as $H_{\text {after }} \downarrow, R_{\text {sequence }} \uparrow$


## Acknowledgements

- Larry Gold
- Gary Stormo
- Andrzej Ehrenfeucht
- Paul Anagnostopoulos


## Version

version $=1.20$ of evtalk.tex 2012 Mar 15


[^0]:    ${ }^{1}$ T. D. Schneider, G. D. Stormo, L. Gold, and A. Ehrenfeucht. J. Mol. Biol., 188:415-431, 1986.
    ${ }^{2}$ R. M. Stephens and T. D. Schneider. J. Mol. Biol., 228:1124-1136, 1992.
    ${ }^{3}$ F. E. Penotti. J Mol Biol, 213:37-52, 1990.

